# Project 5

# **Traveling Salesperson**

### 1. Code included at the end to improve readability

	2. Complexities
Greed	y Algorithm
	Overall Time Complexity: O(n^2)
	Overall Space Complexity: O(n)
	BSSF Initialization
	Time: Constant
	Space: O(n) - stores each city
Brancl	h-and-Bound Algorithm
	<b>Overall Time Complexity:</b> O(2 <sup>n</sup> *n <sup>2</sup> ) for worst case, but closer to O(n <sup>2</sup> * k), where k is the total number of states checked, on average
	Overall Space Complexity: O(n^2)
	Priority Queue - I used the provided heapq Python function
	Time: O(log(n)) for both heappush and heappop as it reheapifies cities
	Space: O(n)
	BSSF Initialization - This is equivalent to the greedy algorithm's runtime.
	Time: O(n^2)
	Space: O(n)
	Reduced Cost Matrix
	Time: O(n^2)
	Space: $O(n^{2})$

### 3. Data Structures for States

I utilized an increasingly incremented integer value to help me keep a tally of all created states as I went, with a separate variable that I also incremented for pruned states. I used a reduced matrix (implemented using numpy) to explore relative costs and see if there are more nodes to explore/prune. This took (n^2) time and space.

# 4. Priority Queue Data Structure

I used the provided heapq package for my priority queue instantiation. Heapq is a binary heap package that allows us to use push and pop functions that will re-heapify to provide us with priority queue capabilities.

## 5. Initial BSSF

For my branch-and-bound algorithm, I used the result of my greedy algorithm to select my initial best solution so far. For relatively small city numbers, our greedy algorithm will run very quickly, and provides us a much better starting point than a random solution might.

Scenario			Greedy		Branch-and-Bound					
Numher of Cities	Seed	Difficulty	Time	Length	Runtime	Cost Best Tour *Optimal	Max # Stored States	# BSSF Updates	# Total Staes	# Total Pruned States
15	20	Hard	0.007	11,081	10.80	9,836*	44	21	15,785	12,829
16	902	Hard	0.007	11,654	58.23	8,051*	60	48	73,557	61,450
18	34	Hard	0.011	10,992	60	9,774	86	8	69,203	56,422
17	1	Hard	0.006	11,747	60	11,636	47	3	76,234	64,979
22	1410	Hard	0.016	13,106	60	15,181	97	0	60,546	52,754
11	1969	Hard	0.004	8,920	1.27	9,110*	19	14	869	638
14	24	Hard	0.007	9,998	6.69	8,716*	31	10	9,922	8,270
18	8	Hard	0.008	12,833	60	12,430	83	6	72,807	60,142
19	19	Hard	0.008	10,730	60	10,301	77	7	67,339	57,674
12	9	Hard	0.002	11,290	2.23	8,649*	33	29	3,660	2,766

### 6. Results Table

# 7. Analysis of Report Table

We seem to observe a factorial increase in states as we increase our city allotment, which is understandable with our worst case time complexity. Total states increase drastically as we increase the city allotment as is expected. We are able to find different solutions and update our BSSF with small city sizes, and on some large counties we were not able to update at all. One fascinating take-away for me is still the overall variability in the data based on the number of cities. This is likely just due to random chance and the fact that our initial state provided by our greedy algorithm is so important, but it is still surprising to see relatively small total states for our 22-city scenario compared to other lesser city count scenarios.

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		CS 3	312 - Winter 2020
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9		Ру	dis. dicore import diiner, drointr
ש. 1			
2		rt i	time
4	impo	rt	
	from	TS	PClasses import *
.6		rt	
	impo	rt	itertools
.8			
		co	by import copy, deepcopy
	clas	s T	SPSolver:
		def	init( self, gui_view ):
3			<pre>selfscenario = None</pre>
4			
5		def	setupWithScenario( self, scenario ):
6			<pre>setfscenario = scenario</pre>
o o			<summary> This is the entry point for the default celver</summary>
.9 10			This is the entry point for the default solver
1			initial RSE
2			
3			<pre>systemations/ eventures/results dictionary for GUI that contains three ints: cost of solution.</pre>
4			time spent to find solution, number of permutations tried during search, the
			solution found, and three null values for fields not used for this
			algorithm
		def	defaultRandomTour( <i>self</i> , <i>time_allowance</i> =60.0 ):
0			results = {}
1			cities = selfscenario.getCities()
2			ncities = [en(cities)
			Toundlour = False
4 E			count = 0
6			ussi - Nunc
7			while not foundfour and time.time()-start time < time allowance:
8			# create a random permutation
.9			perm = np.random.permutation(ncities)
0			route = []
			# Now build the route using the random permutation
2			for i in range(ncities):
			<pre>route.append(cities[perm[i]])</pre>
			bssf = TSPSolution(route)
			count += 1

55		count += 1
56		if bssf.cost < np.inf:
57		# Found a valid route
58		found our = True
50		
60		and time $-$ time $()$
00		end_time = time, time()
01		results[ cost ] = bss1.cost II foundfour else math.in
62		results (time) = end_time - start_time
63		results['count'] = count
64		results['soln'] = bssf
65		results['max'] = None
66		results['total'] = None
67		results['pruned'] = None
68		return results
69		
70		<summarv></summarv>
71		This is the entry point for the greedy solver, which you must implement for
72		the group project (but it is probably a good idea to just do it for the branch-and
73		bound project as a way to get your feet wet). Note this could be used to find your
7/		bound project us a way to get your rect weth. Note this could be used to rind your
75		
75		<pre>&gt;/summary</pre>
70		Are units results difficultating for Gol that contains three ints, cost of best solution,
77		time spent to tind best solution, total number of solutions found, the best
78		solution found, and three null values for fields not used for this
79		algorithm
80		
81		
82		
83		This is a greedy algorithmic approach to finding the solution to the TSP.
84		
85		
86		Time Complexity: $0(n^2)$ – We go through each node and also check all the neightbors of the
87		node to choose our best cost next city. This is $O(n^2)$ and is not optimal
88		Space Complexity: O(n) - We store each city
89		
90		
Q1	def	greedy(self_time_allowance=60_0).
02	uci	Sect - None
92		usia – Noie colf comparia act(ition()
93		cities = Set(_stein) by getettes()
94		notices = ten(cities)
95		startcltypict = {}
96		<pre>start_time = time.time()</pre>
97		
98		# Runs until a solution is found or we go over the allotted time
99		<pre>while (time.time() - start_time) &lt; time_allowance:</pre>
100		
101		
102		
103		for node in range(len(cities)):
104		
105		city = cities[node]
106		cityRoute = []
107		cityRoute.append(city)
108		toVisitCities = deepcopy(cities)
4.00		

108	<pre>toVisitCities = deepcopy(cities)</pre>
109	currCity = city
110	
111	
112	<pre>del toVisitCities[node]</pre>
113	<pre>while len(toVisitCities):</pre>
114	cityCosts = <i>self</i> .closestCities(currCity, toVisitCities)
115	<pre>closestCity = cityCosts[0]</pre>
116	<pre>locClosestCity = toVisitCities.index(closestCity[0])</pre>
117	citvNext = toVisitCities[locClosestCitv]
118	
119	
120	cityRoute.append(cityNext)
121	currCity = cityNext
122	
123	# Removes the next city from our array to visit
124	del toVisitCities[locClosestCity]
125	
126	<pre>if len(toVisitCities):</pre>
127	continue
128	else
129	bssf = TSPSolution(cityRoute)
130	<pre>end time = time.time()</pre>
131	results = {}
132	results['cost'] = bssf.cost
133	results['time'] = end time - start time
134	results['count'] = None
135	results['soln'] = bssf
136	results['max'] = None
137	results['total'] = None
138	results['pruned'] = None
139	<pre>startCityDict[node] = results</pre>
140	continue
141	
142	<pre>self.lowestCost = float("inf")</pre>
143	<pre>for key. solution in startCityDict.items();</pre>
144	<pre>if solution["cost"] &lt; self.lowestCost:</pre>
145	<pre>self.lowestCost = solution["cost"]</pre>
146	lowest = solution
147	
148	return lowest
149	
150	''' <summary></summary>
151	This is the entry point for the branch-and-bound algorithm that you will implement
152	
153	<returns>results dictionary for GUI that contains three ints: cost of best solution,</returns>
154	time spent to find best solution, total number solutions found during search (does
155	
156	
157	
158	

159		
160		Time Complexity: $0(n^2 * 2^n)$ for worst case, but on average closer to $0(k * n^2)$ , where k is the number of states checked
161		Space Complexity: 0(n^2)
162		
163		I also implemented thelt operator on the City class that compares indices
164		to make comparisons easier. It is constant time.
165		
166		
167	def	<pre>branchAndBound(self. time allowance=60.0):</pre>
168		
169		# We are going to use our greedy algorithm to find our initial BSSF
170		# This will give us a better starting point than default
171		<pre>bssf = self.greedv(time allowance=time allowance)['soln']</pre>
172		
173		cities = self, scenario $def(ities())$
174		cities = cities
175		
176		ncities - len(cities)
177		hear = 1
178		hcap = []
170		$b_{s}(b_{s}) = 0$
100		
101		
101		
102		humotottons - V
107		
104		# Cost for initializing reduced matrix is $O(n^2)$
196		inited with a second seco
107		Intreducednati IX, tower bound = Set). Intra (Reducednati IX(CITES)
199		starting = tup/a/lowerRound_cities[0]_cities[1]_initPeducedMatrix [cities[0]_indev]_lowerRound)
180		Start Ling = Laber (ower bound, Cittes[0], Cittes[1,], Hitkeducedmatrix, [Cittes[0], _index], (ower bound)
100		time ctart - time time ()
101		
102		# We will run until our bean is empty, or we go overe our time-allotment (60 seconds by default)
192		while ((time time () - time start) < time allowance and len(hean)).
194		
195		# Time complexity for beand's beannon is $\Omega(\log(n))$
196		$n \operatorname{Derk} = \operatorname{heard}_{n} \operatorname{heard}_{n} (\operatorname{heard}_{n})$
197		ondeex = neuppineupponeup)
198		if onDeck[5] < self.lowestCost:
199		
200		# We will iterate through each city we still need to visit
201		for city in onDeck[2].
202		
202		# We need to check if the path event exits (has an edge)
205		if salf, scencio, educe viet (onleck [1], index][city, index].
205		
206		# Creation of our reduced matrix is $\Omega(n^2)$
207		newSubtree = self.reducedMatrix(rity, onDeck[3], onDeck)
208		
209		# checks to see if we have more cities to visit
210		if not len(newSubtree[2]):
211		
_		



263		
264	def	<pre>convertIndicesToCities(self, indicesCities):</pre>
265		cityList = []
266		for index in indicesCities:
267		citylist_append(self_cities[index])
269		
200		
209		
270		
271		
272		
273		Time Complexity: O(n)
274		Space Complexity: 0(1)
275		
276		
277	def	clearCity(self. toVisitCities. toVisitNext):
278		for index. city in enumerate(toVisitCities):
270		if city index — talisitaevi index:
280		
200		
201		Dieak
282		
283		det tovisituities[clearindex]
284		return tovisitCitles
285		
286		
287		Time Complexity: O(n)
288		Space Complexity: O(n)
289		
290		
291	def	closestCities(self, city, cityList):
292		$cost = \{\}$
293		for thy sit ( it is in cityl ist.
200		cont[toVisitfities] = city_costTo(toVisitfities)
294		
295		$b_{1}$
290		by cost (10) and the cost (10) (a cost (10) (a cost (10))
297		<pre># print( closest length is {}".format(bycostcities[0][i]))</pre>
298		return bycostcitles
299		
300		
301		Time Complexity: O(n^2) to compare distance of node from every other ndoe
302		Space Complexity: 0(n^2) to store all relative distances
303		
304		
305	def	reducedMatrix(self. toVisitNext. matrix. scenario):
306		subtree = deepcopy(scenario)
307		matrix = matrix.copy()
308		
300		initCost = matrix[cubtree[1] index][to//icitNevt index]
210		
310		
212		
312		matrix[subtree[1]index] = np.int
313		matrix[:, tovisitNextindex] = np.int
314		<pre>matrix[toVisitNextindex][subtree[1]index] = np.inf</pre>
315		
D 0 0		

315		
316		
317		# Reduce columns and rows
318		<pre>for row in range(matrix.shape[0]):</pre>
319		rowMin = np.min(matrix[row])
320		if np.isinf(rowMin):
321		
322		matrix[row] = matrix[row] - rowMin
323		reduceSum += rowMin
324		
325		
326		<pre>for col in range(matrix.shape[1]):</pre>
327		colMin = np.min(matrix[:, col])
328		if np.isinf(colMin):
329		continue
330		matrix[; col] = matrix[; col] - coumin
331		reaucesum += coumin
332 333		
227		tovisite(tites = subtree[z])
334		tovisitettes - set clearcity(tovisitettes, tovisitext)
336		<pre>newCost = subtree[5] + initCost + reduceSum</pre>
337		
338		# Value for the queue
339		newValue = score / len(cities_visited)
340		
341		
342		distanceTuple = {newValue, toVisitNext, toVisitCities, matrix, subtree[4] + [toVisitNextindex], newCost)
343		
344		
345		
347		Time Complexity: 0(n^2) to compare distance of node from every other ndoe
348		Space Complexity: 0(n <sup>2</sup> ) to store all relative distances
349		
350		
351	def	initalReducedMatrix( <i>self, cityList</i> ):
352		
353		# We will initialize our matrix, with infinity
354		<pre>matrix = np.full((len(cityList), len(cityList)), fill_value=np.inf)</pre>
355		
350		# We will tind relative distances, all diagnots will remain intinity
358		for from Loc, city in enumerate(cityList):
350		for toloc, toCity in enumerate(citylist):
360		
361		if fromLoc == toLoc:
362		
363		
364		
365		<pre>matrix[tromLoc][toLoc] = city.costlo(toCity)</pre>
360		
507		
367		
368		reducesum = 0
369		# do rows and get mins
370		row row and range(matrix.shape(0)):
371		matrix/fowl_n = matrix/fowl_n = couMin
372		
374		
375		# get reduced columns now
376		for col in range(matrix.shape[1]):
377		colMin = np.min(matrix[:, col])
378		<pre>matrix[:, col] = matrix[:, col] - colMin</pre>
379		reduceSum += colMin
380		
380 381		return matrix, reduceSum
380 381 382 383		return matrix, reduceSum
380 381 382 383 384		return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project.</summary>
380 381 382 383 384 385		return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project. </summary>
380 381 382 383 384 385 386		return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project. </summary> <returns>results dictionary for GUI that contains three <u>ints: cost of best solution</u>,</returns>
380 381 382 383 384 385 386 387		return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project. </summary> <returns>results dictionary for GUI that contains three ints: cost of best solution, time spent to find best solution, total number of solutions found during search, the</returns>
380 381 382 383 384 385 386 386 387 388		return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project. </summary> <returns>results dictionary for GUI that contains three ints: cost of best solution, time spent to find best solution, total number of solutions found during search, the best solution found. You may use the other three field however you like.</returns>
380 381 382 383 384 385 386 387 388 388 389		return matrix, reduceSum <pre>     summary&gt;     This is the entry point for the algorithm you'll write for your group project.      <returns>results dictionary for GUI that contains three ints: cost of best solution,     time spent to find best solution, total number of solutions found during search, the     best solution found. You may use the other three field however you like.     algorithm</returns> </pre>
380 381 382 383 384 385 386 387 388 389 390		return matrix, reduceSum <pre>     summary&gt;     This is the entry point for the algorithm you'll write for your group project.      <returns>results dictionary for GUI that contains three ints: cost of best solution,     time spent to find best solution, total number of solutions found during search, the     best solution found. You may use the other three field however you like.     algorithm</returns></pre>
380 381 382 383 384 385 386 387 388 389 390 391		<pre>return matrix, reduceSum  <pre> summary&gt; This is the entry point for the algorithm you'll write for your group project.  <returns>results dictionary for GUI that contains three ints: cost of best solution, time spent to find best solution, total number of solutions found during search, the best solution found. You may use the other three field however you like. algorithm</returns> Because of complications with Covid-19, we are not going to complete this </pre></pre>
380 381 382 383 384 385 386 387 388 387 388 389 390 391 392		<pre>return matrix, reduceSum  <pre>csummary&gt; This is the entry point for the algorithm you'll write for your group project.  <returns>results dictionary for GUI that contains three ints: cost of best solution, time spent to find best solution, total number of solutions found during search, the best solution found. You may use the other three field however you like. algorithm</returns> Because of complications with Covid-19, we are not going to complete this portion of the project</pre></pre>
380 381 382 383 384 385 386 387 388 387 388 389 390 391 391 392 393 394		<pre>return matrix, reduceSum ' <summary> This is the entry point for the algorithm you'll write for your group project. </summary> <returns>results dictionary for GUI that contains three ints: cost of best solution, time spent to find best solution, total number of solutions found during search, the best solution found. You may use the other three field however you like. algorithm</returns> Because of complications with Covid-19, we are not going to complete this portion of the project f fancy(self_time_allowance=60.0). </pre>

def fancy(self, time\_allowance=60.0):
 print("Not yet implemented")